

Information Theory w/ Applications

Fall 2008 - Homework 1

1. Let $Z: \Omega \rightarrow \mathbb{R}$, then for $\tau > 0$, $a \in \mathbb{R}$

$$\mathbb{P}(Z \geq a) = \mathbb{P}(e^{\tau Z} \geq e^{\tau a})$$

$$= \mathbb{E} \left[\mathbb{1}_{e^{\tau Z} \geq e^{\tau a}} \right]$$

$$\leq \mathbb{E} \left[\frac{e^{\tau Z}}{e^{\tau a}} \right] = e^{-\tau a} \mathbb{E} [e^{\tau Z}].$$

$$\begin{aligned} 2. a) H(X, Y, Z) &= H(X) + H(Y|X) + H(Z|X, Y) \\ &= H(X) + H(Z|X) + H(Y|X, Z) \\ &= H(Y) + H(X|Y) + H(Z|X, Y) \\ &= H(Y) + H(Z|Y) + H(X|Y, Z) \\ &= H(Z) + H(X|Z) + H(Y|X, Z) \\ &= H(Z) + H(Y|Z) + H(X|Y, Z) \end{aligned}$$

$$b) \quad H(X) = H(P_X)$$

$$\text{where } P_X(0) = P_X(1) = \frac{1}{2}$$

$$\Rightarrow H(X) = h\left(\frac{1}{2}\right) = 1 \text{ bit}$$

$$H(Y|X=1) = 0 \quad \text{b/c } P_Y(0|X=1) = 1$$

$$H(Y|X=0) = 1 \text{ bit} \quad \text{b/c } P_Y(0|X=0) = \frac{1}{2}$$

so

$$H(Y|X) = \frac{1}{2}(0) + \frac{1}{2}(1) \text{ bit}$$

~~H~~ Now

$$P(Z|X=a, Y=b) = \begin{cases} 1, & a=b=0 \text{ or } b=1 \\ \frac{1}{2}, & a=1, b=0 \\ 0, & \text{else} \end{cases}$$

gives

$$H(Z|X=0, Y=0) = 0$$

$$H(Z|X=0, Y=1) = 0$$

$$H(Z|X=1, Y=0) = 1 \text{ bit}$$

so

$$H(Z|X, Y) = \frac{1}{2}(1) \text{ bit}$$

c) We conclude

$$\begin{aligned} H(X, Y, Z) &= 1 + \frac{1}{2} + \frac{1}{2} \text{ a bit} \\ &= 2 \text{ bit,} \end{aligned}$$

which is (joint) entropy of prob.
measure with 4 equal-prob.
outcomes

$$H_4 = \sum_{j=1}^4 \left(\frac{1}{4}\right) \log_2\left(\frac{1}{4}\right) = 2 \text{ bit.}$$

$$3. H(X|Z) = H(X, Y|Z) - H(Y|X, Z)$$

$$\leq H(X, Y|Z)$$

$$= H(X|Y, Z) + H(Y|Z)$$

$$\leq H(X|Y) + H(Y|Z)$$

$$4. \quad a) \quad C_0(P, Q) = -\log \underbrace{\sum_{j=1}^m q_j}_1 = 0$$

$$C_1(P, Q) = 0 \quad \text{by symmetry}$$

$$C_\alpha(P, P) = -\log \underbrace{\sum_{j=1}^m P_j}_1 = 0$$

b) By Hölder's ineq.,

$$\text{setting } p\alpha = q(1-\alpha) = 1 \Rightarrow \frac{1}{p} + \frac{1}{q} = 1$$

gives

$$\begin{aligned} & \sum_{j=1}^m P_j^\alpha q_j^{1-\alpha} \\ & \leq \left(\underbrace{\sum_{j=1}^m P_j}_{1}^{\frac{p\alpha}{\alpha}} \right)^{\frac{1}{p}} \left(\underbrace{\sum_{j=1}^m q_j}_{1}^{\frac{q(1-\alpha)}{1-\alpha}} \right)^{\frac{1}{q}} = 1 \end{aligned}$$

so

$$C_\alpha(P, Q) \geq 0 .$$

$$c) \frac{\partial}{\partial \alpha} C_\alpha(P, Q) = - \frac{1}{\sum_{i=1}^n P_i^\alpha q_i^{1-\alpha}} \left(\sum_{i=1}^n (P_i^\alpha (\log P_i) q_i^{1-\alpha} - P_i^\alpha \log\left(\frac{q_i}{P_i}\right) q_i^{1-\alpha}) \right)$$

$$= - (\log \alpha + \log(1-\alpha))$$

$$= \mathbb{E}_\alpha [-(\log P_i + \log q_i)]$$

$$\frac{\partial^2}{\partial \alpha^2} C_\alpha(P, Q) = \frac{1}{\left(\sum_{i=1}^n P_i^\alpha q_i^{1-\alpha} \right)^2} \left(\sum_{i=1}^n (P_i^\alpha \log P_i q_i^{1-\alpha} - P_i^\alpha \log q_i q_i^{1-\alpha})^2 \right)$$

$$- \frac{1}{\sum_{i=1}^n P_i^\alpha q_i^{1-\alpha}} \sum_{i=1}^n \left(P_i^\alpha (\log P_i)^2 q_i^{1-\alpha} - 2 P_i^\alpha \log P_i \log q_i q_i^{1-\alpha} + P_i^\alpha (\log q_i)^2 q_i^{1-\alpha} \right)$$

By Cauchy - Schwarz,

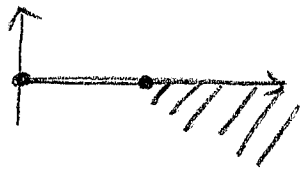
$$\langle \log P_i, \log q_i \rangle_\alpha \leq \| \log P_i \|_\alpha \| \log q_i \|_\alpha \leq \frac{1}{2} \| \log P_i \|_\alpha^2 + \frac{1}{2} \| \log q_i \|_\alpha^2$$

$$\text{where } \langle \cdot, \cdot \rangle = \sum_{i=1}^n P_i^\alpha q_i^{1-\alpha} (\cdot \cdot)$$

$$\Rightarrow \sum_{i=1}^n (\dots) \geq 0 \Rightarrow \frac{\partial^2}{\partial \alpha^2} C_\alpha(P, Q) \leq 0$$

$\Rightarrow C_\alpha(P, Q)$ concave.

d) By concavity and $C_0(P, Q) = C_1(P, Q)$,



$$C_\alpha(P, Q) \leq 0 \quad \text{for} \\ \alpha \geq 1$$

e) By ~~concavity~~ Hölder, if $P \neq Q$ and $\alpha \notin \{0, 1\}$ $0 < \alpha < 1$, then $C_\alpha(P, Q) > 0$.

But since $C_1(P, Q) = 0$ and

$C_\alpha(P, Q)$ is concave, $C_\alpha(P, Q) < 0$

for $\alpha > 1$. Otherwise, $P = Q$

implies $C_\alpha(P, Q) = 0$.

f) Consider α^* such that $C_{\alpha^*}(P, Q)$ maximal, then $C_{\alpha^*}(P, Q) = D(S_{\alpha^*} \| P) = D(S_{\alpha^*} \| Q)$.

From c) max. is unique and

$$\sum_{i=1}^n P_i^\alpha q_i^{1-\alpha} \underbrace{(\log P_i - \log q_i)}_{\log \frac{P_i}{q_i}} = 0$$

$$\Rightarrow \frac{1}{N} \sum_{i=1}^n P_i^\alpha q_i^{1-\alpha} \log P_i = \frac{1}{N} \sum_{i=1}^n P_i^\alpha q_i^{1-\alpha} \log q_i$$

where N is normalization factor.

Now subtracting ~~from~~ both sides ~~from~~ from

$$\frac{1}{N} \sum_{i=1}^n p_i^\alpha q_i^{1-\alpha} \log \frac{p_i^\alpha q_i^{1-\alpha}}{N},$$

we obtain $D(S_{\alpha^*} \| P) = D(S_{\alpha^*} \| Q)$.

5. Given acceptance region \mathcal{A}'_n

and

$$\alpha'_n = P(X_1 \notin \mathcal{A}'_n)$$

$$\beta'_n = P(X_2 \in \mathcal{A}'_n),$$

we want to minimize the error probability

$$P(\theta \neq \phi(X^n)) = \pi_1 \alpha'_n + \pi_2 \beta'_n.$$

$$= \pi_1 \sum_{x \notin \mathcal{A}'_n} P_1(x) + \pi_2 \underbrace{\sum_{x \in \mathcal{A}'_n} P_2(x)}$$

$$\pi_2 \left(1 - \sum_{x \notin \mathcal{A}'_n} P_2(x) \right)$$

$$= \pi_1 \left(\sum_{x \notin \mathcal{A}'_n} (P_1(x) - \frac{\pi_2}{\pi_1} P_2(x)) \right) + \pi_2$$

so minimizer has acceptance region

$$A_n = \{x \in \mathbb{A}^n : \pi_1 P_1(x) \geq \pi_2 P_2(x)\}$$

For given $x \in \mathbb{A}^n$, let $\mu_x(a) = \frac{1}{n} |\{j : x_j = a\}|$
then

$$\frac{P_1(x)}{P_2(x)} \geq \frac{\pi_2}{\pi_1}$$

is equivalent to componentwise ($x = (x_1, x_2, \dots, x_n)$)

$$\sum_{j=1}^n \log \frac{P_1(x_j)}{P_2(x_j)} \geq \log \frac{\pi_2}{\pi_1}$$

and

$$\sum_{a \in \mathbb{A}} \mu_x(a) \log \frac{P_1(a)}{P_2(a)} \geq \frac{1}{n} \log \frac{\pi_2}{\pi_1} \equiv \alpha_n \rightarrow 0.$$

By law of large numbers, $\mu_x(a) \xrightarrow{n \rightarrow \infty} P_j(a)$ a.s.

where $j = \theta$. Moreover, using the result of

2.31 b) we have

$$- \lim_{n \rightarrow \infty} \frac{1}{n} \log \alpha_n = D(R_0^\alpha \parallel P_1)$$

$$- \lim_{n \rightarrow \infty} \frac{1}{n} \log \beta_n = D(R_0^\alpha \parallel P_2)$$

where $\sum_{a \in \mathbb{A}} R_0^\alpha(a) \log \frac{P_1(a)}{P_2(a)} = 0$ and

$$R_0^\alpha(a) = \frac{P_1^\alpha(a) P_2^{1-\alpha}(a)}{\sum_{a \in \mathbb{A}} P_1^\alpha(a) P_2^{1-\alpha}(a)} \quad \text{for some } \alpha.$$

From

$$\begin{aligned} \mathbb{P}(\theta \neq \phi(X^n)) &= \pi_1 \alpha_n + \pi_2 \beta_n \\ &\leq 2 \max \{ \pi_1 \alpha_n, \pi_2 \beta_n \} \end{aligned}$$

and

$$\mathbb{P}(\theta \neq \phi(X^n)) \geq \max \{ \pi_1 \alpha_n, \pi_2 \beta_n \},$$

the optimal test minimizes for $n \rightarrow \infty$

$$\max \left\{ \frac{1}{n} \log \alpha_n, \frac{1}{n} \log \beta_n \right\}.$$

Consequently, α is chosen to maximize

$$\min \{ D(R_0^\alpha \| P_1), D(R_0^\alpha \| P_2) \}.$$

Since $D(R_0^\alpha \| P_1)$ is monotonically decreasing,
 $D(R_0^\alpha \| P_2)$ is increasing,

this max is achieved when $D(R_0^\alpha \| P_1) = D(R_0^\alpha \| P_2)$.

Now, by 2.28 f, we know then

$$D(R_0^\alpha \| P_1) = \max_{\alpha} C_{\alpha}(P_1, P_2).$$