# Thermodynamics for non-uniformly mixing systems: factors of $\beta$ -shifts are intrinsically ergodic

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November 11, 2010

Joint work with Daniel Thompson

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- Introduction
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  - $\beta$ -shifts
  - Intrinsic ergodicity for factors
- General result
  - Specification and CGC-decompositions
  - A criterion for intrinsic ergodicity that passes to factors
  - Other examples

## Basic thermodynamic concepts

Topological dynamical system:

- X a compact metric space,  $f: X \to X$  continuous
- $\mathcal{M} = \{ \text{Borel } f\text{-invariant probability measures on } X \}$

Variational principle:  $h_{\mathrm{top}}\left(X,f
ight)=\sup_{\mu\in\mathcal{M}}h_{\mu}(f)$ 

- If  $h_{\mu}(f) = h_{\text{top}}(X, f)$ , then  $\mu$  is a measure of maximal entropy (MME)
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When is a transitive dynamical system intrinsically ergodic?



#### Motivation and context

More general variational principle for topological pressure  $P(\varphi)$  of a continuous potential function  $\varphi \colon X \to \mathbb{R}$ 

$$P(\varphi) = \sup_{\mu \in \mathcal{M}} \left( h_{\mu}(f) + \int \varphi \, d\mu \right)$$

If  $h_{\mu}(f) + \int \varphi \, d\mu = P(\varphi)$ , then  $\mu$  is an equilibrium state.

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If  $h_{\mu}(f) + \int \varphi \, d\mu = P(\varphi)$ , then  $\mu$  is an equilibrium state.

- Existence of a unique equilibrium state is connected to statistical properties, large deviations, multifractal analysis, phase transitions, etc.
- $\varphi \equiv 0$ : reduces to intrinsic ergodicity. Techniques for showing intrinsic ergodicity usually generalise to help prove other thermodynamic results.

# Intrinsic ergodicity for shift spaces

Focus on shift spaces (subshifts):

- $X \subset \Sigma_p$  or  $X \subset \Sigma_p^+$ , X closed and  $\sigma$ -invariant
- $\mathcal{L} = \mathcal{L}(X) = \{x_1 \cdots x_n \mid x \in X, n \ge 1\}$  is the language of X

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When is a transitive shift space intrinsically ergodic? Not always.

Example:  $X \subset \Sigma_5 = \{0, 1, 2, 1, 2\}^{\mathbb{Z}}$ . Define the language  $\mathcal{L}$  by  $\mathbf{v}0^n \mathbf{w}, \ \mathbf{w}0^n \mathbf{v} \in \mathcal{L}$  if and only if  $n \geq 2 \max(|\mathbf{v}|, |\mathbf{w}|)$ .

- $(X, \sigma)$  is topologically transitive
- $h_{\text{top}}(X, \sigma) = \log 2$
- 2 measures of maximal entropy:

$$\begin{split} & \textcolor{red}{\nu} = (\frac{1}{2}, \frac{1}{2}) \text{-Bernoulli on } \{\textcolor{red}{1,2}\}^{\mathbb{Z}}, \\ & \mu = (\frac{1}{2}, \frac{1}{2}) \text{-Bernoulli on } \{\textcolor{red}{1,2}\}^{\mathbb{Z}}. \end{split}$$

The following classes of shift spaces are intrinsically ergodic:

• Irreducible subshifts of finite type (Parry 1964)



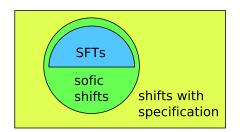
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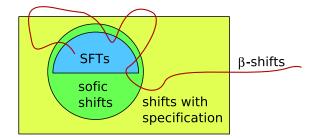
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- $\beta$ -shifts (Walters 1978, Hofbauer 1979)

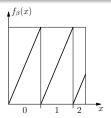


### $\beta$ -shifts

 $\beta>1$ ,  $b=\lceil\beta\rceil$ . The  $\beta$ -shift  $\Sigma_{\beta}\subset\Sigma_{b}^{+}$  is the natural coding space for the map

$$f_{\beta} \colon [0,1] \to [0,1], \qquad x \mapsto \beta x \pmod{1}$$

$$1_{eta} = a_1 a_2 \cdots$$
 , where  $1 = \sum_{n=1}^{\infty} a_n eta^{-n}$ 

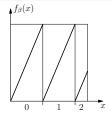


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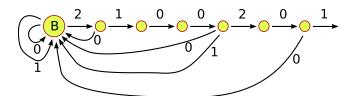
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Fact: Sequences  $x \in \Sigma_{\beta}$  are precisely those sequences in  $\Sigma_b$  that label trajectories of the following graph beginning at the vertex **B**. (Here  $1_{\beta} = 2100201...$ )



Intrinsic ergodicity is not necessarily preserved by factors.

- $X \subset \{0, 1, 2, 1, 2\}^{\mathbb{Z}}$  as before
- $\bullet \ Y \subset \Sigma_6 = \{0, \textcolor{red}{1}, \textcolor{red}{2}, \textcolor{black}{1}, \textcolor{black}{2}, \textcolor{black}{3}\}^{\mathbb{Z}}$  by similar rule
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What intrinsically ergodic classes are closed under factors?

- Closure of SFTs is class of sofic systems
- Specification preserved by factors
- Factors of  $\beta$ -shifts = ?????

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#### Theorem (C.-Thompson 2010)

Yes.



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Are factors of  $\beta$ -shifts intrinsically ergodic? (Klaus Thomsen)

#### Theorem (C.-Thompson 2010)

Every subshift factor of a  $\beta$ -shift is intrinsically ergodic.



# The classical specification property

- $\mathcal{L} = \text{language for a shift space } X$
- $\mathcal{L} \leftrightarrow \{\text{cylinders in } X\}$
- |w| = length of w,  $\mathcal{L}_n = \{ w \in \mathcal{L} \mid |w| = n \}$

X has specification if there exists  $t \in \mathbb{N}$  such that for every  $w_1, \ldots, w_m \in \mathcal{L}$ , there exist  $z_1, \ldots, z_{m-1} \in \mathcal{L}_t$  for which the concatenated word  $w_1 z_1 w_2 z_2 \cdots z_{m-1} w_m$  is in  $\mathcal{L}$ .

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Topological transitivity guarantees the existence of such words  $z_i \in \mathcal{L}$ . Specification demands that the words  $z_i$  can be chosen to have uniformly bounded length t, where t is independent of the words  $w_i$  and their lengths.

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 $\Sigma_{\beta}$  does not have the specification property if  $1_{\beta}$  contains arbitrarily long strings of 0's.



The following shifts have the specification property:

- Mixing subshifts of finite type
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- Some β-shifts

 $\Sigma_{\beta}$  does not have the specification property if  $1_{\beta}$  contains arbitrarily long strings of 0's.



 $\Sigma_{\beta}$  does not have specification for Lebesgue-a.e.  $\beta > 1$ .

We must replace specification with a property that

- holds for every  $\beta$ -shift;
- implies intrinsic ergodicity;
- is preserved by factors.

# A restricted version of the specification property

Fix a subset  $\mathcal{G} \subset \mathcal{L}$ . We say that  $\mathcal{G}$  has specification if there exists  $t \in \mathbb{N}$  such that for every  $w_1, \ldots, w_m \in \mathcal{G}$ , there exist  $z_1, \ldots, z_{m-1} \in \mathcal{L}_t$  for which the concatenated word  $x := w_1 z_1 w_2 z_2 \cdots z_{m-1} w_m$  is in  $\mathcal{L}$ .

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Example: For  $X = \Sigma_{\beta}$ , let  $\mathcal{G}$  be the set of words corresponding to paths that begin and end at **B**. Then  $\mathcal{G}$  has (Per)-specification with t = 0.

A CGC-decomposition of the language  $\mathcal{L}$  is a collection of words  $\mathcal{C}^p, \mathcal{G}, \mathcal{C}^s \subset \mathcal{L}$  with the following properties.

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Example: For  $X = \Sigma_{\beta}$ , let  $\mathcal{C}^p = \emptyset$  and let  $\mathcal{C}^s$  be the set of words corresponding to paths that begin at **B** and never return. Then  $(\mathcal{C}^p, \mathcal{G}, \mathcal{C}^s)$  is a uniform CGC-decomposition.

# Intrinsic ergodicity for shifts with CGC-decompositions

Given a collection of words  $\mathcal{D} \subset \mathcal{L}$ , let  $h(\mathcal{D}) = \overline{\lim}_{n \to \infty} \frac{1}{n} \log \# \mathcal{D}_n$ . Observe that  $h_{\text{top}}(X, \sigma) = h(\mathcal{L})$ .

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#### Theorem (C.-Thompson 2010)

Let X be a shift space admitting a uniform CGC-decomposition. If  $h(\mathcal{C}^p \cup \mathcal{C}^s) < h_{\mathrm{top}}(X,\sigma)$ , then  $(X,\sigma)$  is intrinsically ergodic. If  $\mathcal{G}$  has (Per)-specification, then the unique MME is the limit of the periodic orbit measures  $\mu_n = \frac{1}{\#\{x|f^n(x)=x\}} \sum_{f^n(x)=x} \delta_x$ .

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Example: For  $X = \Sigma_{\beta}$ , let  $x = 1_{\beta}$ . Then  $(\mathcal{C}^p \cup \mathcal{C}^s)_n = \{x_1 \cdots x_n\}$ , and so  $h(\mathcal{C}^p \cup \mathcal{C}^s) = 0$ . Thus  $(\Sigma_{\beta}, \sigma)$  is intrinsically ergodic.

#### Behaviour under factors

Let  $(\tilde{X}, \sigma)$  be a factor of  $(X, \sigma)$ , and let  $\mathcal{L}, \tilde{\mathcal{L}}$  be the languages.

• If  $\mathcal{L}$  has a uniform CGC-decomposition, then so does  $\tilde{\mathcal{L}}$ . Futhermore,  $h(\tilde{\mathcal{C}}^p \cup \tilde{\mathcal{C}}^s) \leq h(\mathcal{C}^p \cup \mathcal{C}^s)$ .

Every factor with  $h_{\text{top}}(\tilde{X}, \sigma) > h(\mathcal{C}^p \cup \mathcal{C}^s)$  is intrinsically ergodic.

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Dichotomy for shifts with uniform CGC-decompositions:

Either  $h_{\text{top}}(X, \sigma) > 0$ , or X comprises a single periodic orbit.

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Let X be a shift space admitting a uniform CGC-decomposition.

If  $h(C^p \cup C^s) = 0$ , then every subshift factor of  $(X, \sigma)$  is intrinsically ergodic.

# S-gap shifts

Fix  $S \subset \mathbb{N}$  and suppose S is infinite. The associated S-gap shift is the subshift  $\Sigma_S \subset \{0,1\}^{\mathbb{Z}}$  with language

$$\mathcal{L} = \{0^k 10^{n_1} 10^{n_2} 1 \cdots 10^{n_j} 10^{\ell} \mid n_i \in S, k, \ell \in \mathbb{N}\}.$$

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A uniform CGC-decomposition for  $\Sigma_S$  is given by

$$\mathcal{G} = \{0^n 1 \mid n \in S\}$$

$$\mathcal{C}^p = \{0^k 1 \mid k \ge 0\}$$

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Then  $\#(\mathcal{C}^p \cup \mathcal{C}^s)_n = 2$  for all  $n \ge 1$ , and so  $h(\mathcal{C}^p \cup \mathcal{C}^s) = 0$ . It follows that every subshift factor of an S-gap shift is intrinsically ergodic.

## Coded systems

A shift space X is coded if its language  $\mathcal{L}$  is freely generated by a countable set of generators  $\{w_n\}_{n\in\mathbb{N}}\subset\mathcal{L}$ .

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Every coded system has a uniform CGC-decomposition.

$$\mathcal{G} = \{ w_{n_1} w_{n_2} \cdots w_{n_k} \mid n_i \in \mathbb{N} \}$$

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# Coded systems

A shift space X is coded if its language  $\mathcal{L}$  is freely generated by a countable set of generators  $\{w_n\}_{n\in\mathbb{N}}\subset\mathcal{L}$ .

$$\mathcal{L} = \{ \text{all subwords of } w_{n_1} w_{n_2} \cdots w_{n_k} \mid n_i \in \mathbb{N} \}$$

Every coded system has a uniform CGC-decomposition.

$$\mathcal{G} = \{ w_{n_1} w_{n_2} \cdots w_{n_k} \mid n_i \in \mathbb{N} \}$$

$$\mathcal{C}^p = \{ \text{suffixes of } w_n \mid n \in \mathbb{N} \}$$

$$\mathcal{C}^s = \{ \text{prefixes of } w_n \mid n \in \mathbb{N} \}$$

Let  $\hat{h} = h(\{\text{prefixes and suffixes of generators}\})$ .

- $\hat{h} < h_{\text{top}}(X, \sigma) \Rightarrow (X, \sigma)$  is intrinsically ergodic
- $\hat{h} = 0 \Rightarrow$  every subshift factor of  $(X, \sigma)$  is intrinsically ergodic